

# 1 Breaking of Simplified Data Encryption Standard Using Genetic 2 Algorithm

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## 7 Abstract

8 Cryptanalysis of ciphertext by using evolutionary algorithm has gained so much interest in  
9 recent years. In this paper we have used a Genetic algorithm with improved crossover  
10 operator (Ring Crossover) for cryptanalysis of SDES. There so many attacks in cryptography.  
11 The cipher text attack only is considered here and several keys are generated in the different  
12 run of the genetic algorithm on the basis of their cost function value which depends upon  
13 frequency of the letters. The results on the S-DES indicate that, this is a promising method  
14 and can be adopted to handle other complex block ciphers like DES, AES.

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16 **Index terms**— Cryptanalysis, Ciphertext attack, Simplified Data Encryption Standard, genetic algorithm,  
17 Key search space

18 cipher is a secret way of writing in which plaintext is converted into a scrambled (encrypted) version of the  
19 original message (ciphertext) by using a key. Those who know the key can easily decrypt the ciphertext back into  
20 the plaintext. Cryptanalysis is the study of breaking ciphers that is finding the key or converting the ciphertext  
21 into the plaintext without knowing the key. Many cryptographic systems have a finite key space and, hence, are  
22 vulnerable to an exhaustive key search attack. Yet, these systems remain secure from such an attack because the  
23 size of the key space is such that the time and resources for a search are not available. Optimization techniques  
24 have got a significant importance in determining efficient solutions of different complex problems. One such  
25 problem is to break S-DES. This paper considers cryptanalysis of S-DES. In the brute force attack, the attacker  
26 tries each and every possible key on the part of cipher text until desired plaintext is obtained. A brute force  
27 approach may take so much time to guess the real key which is used to generate a cipher text, so the difficulty  
28 of breaking the cipher is directly proportional to the number of keys. On the other hand optimization technique  
29 can be used for the same purpose. Genetic algorithm is an evolutionary algorithm that works well and takes less  
30 time to break cipher as compared to Brute force attack.

31 The remaining paper is organized as follows: Section II discusses the earlier works done in this field. Section  
32 III presents overview of S-DES and Section IV gives the overview of Genetic Algorithm. Experimental results  
33 are discussed in Section V. Conclusion are presented in section VI At last References are given.

34 In the last few years, so many papers have been published in the field of cryptanalysis. R.Spillman etc.  
35 showed that Knapsack cipher [4] and substitution ciphers [5] could be attacked using genetic algorithm. In  
36 the recent years Garg[1,2] presented the use of memetic algorithm and genetic algorithm to break a simplified  
37 data encryption standard algorithm. Nalini [3] used efficient heuristics to attack S-DES. In 2006 Nalini used  
38 GA, Tabu search and Simulated Annealing techniques to break S-DES. Matusi [7] showed the first experimental  
39 cryptanalysis of DES using an linear cryptanalysis technique. Clark [6] also presented important analysis on how  
40 different optimization techniques can be used in the field of cryptanalysis. Vimalathithan [9] also used GA to  
41 attack Simplified-DES. In this paper, a Genetic Algorithm with improved parameters is used to break S-DES. A  
42 population of keys is generated and their fitness is calculated by using efficient fitness function. At the end, we  
43 will find the key in less time.

44 In this section we will provide the overview of S-DES Algorithm. Simplified DES, developed by Professor  
45 Edward Schaefer of Santa Clara University is an educational rather than a secure encryption algorithm. The

46 S-DES [8,10] encryption algorithm takes an 8-bit block of plaintext and a 10-bit key as input and produces an  
 47 8-bit block of ciphertext as output. The S-DES decryption algorithm takes an 8-bit block of ciphertext and  
 48 the same 10-bit key used to produce that ciphertext as input and produces the original 8-bit block of plaintext.  
 49 The encryption algorithm involves five functions: an initial permutation (IP); a complex function labeled  $f_K$ ,  
 50 which involves both permutation and substitution operations and depends on a key input; a simple permutation  
 51 function that switches (SW) the two halves of the data; the function  $f_K$  again; and finally a permutation function  
 52 that is the inverse of the initial permutation (IP<sup>-1</sup>).

53 The function  $f_K$  takes as input not only the data passing through the encryption algorithm, but also an 8-bit  
 54 key. S-DES uses a 10-bit key from which two 8-bit subkeys are generated. In this, the key is first subjected to a  
 55 permutation (P10). Then a shift operation is performed. The output of the shift operation then passes through  
 56 a permutation function that produces an 8-bit output (P8) for the first subkey ( $K_1$ ). The output of the shift  
 57 operation also feeds into another shift and another instance of P8 to produce the second subkey ( $K_2$ ). The  
 58 S-boxes operate as follows:

59 The first and fourth input bits are treated as 2-bit numbers that specify a row of the S-box and the second and  
 60 third input bits specify a column of the S-box. The entry in that row and column in base 2 is the 2-bit output.

## 1 The Switch Function

62 The function  $f_K$  only alters the leftmost 4 bits of the input. The switch function (SW) interchanges the left and  
 63 right 4 bits so that the second instance of  $f_K$  operates on a different 4 bits. In this second instance, the E/P,  
 64 S0, S1, and P4 functions are the same. The key input is  $K_2$ .

65 The genetic algorithm [13,20] is a search algorithm based on the natural selection and on "survival of the  
 66 fittest", the main idea is that in order for a population of individuals to adapt to some environment, it should  
 67 behave like a natural system. This means that survival and reproduction of an individual is promoted by the  
 68 elimination of useless traits and by rewarding useful behavior. The genetic algorithm belongs to the family of  
 69 evolutionary algorithms. An evolutionary algorithm maintains a population of solutions for the problem at hand.  
 70 The population is then evolved by the iterative application of a set of stochastic operators. The simplest form  
 71 of genetic algorithm involves three types of operators: selection, crossover and mutation. A selection operator is  
 72 applied first. In this paper, we are using Ring crossover operator [11]. In ring crossover two parents such as parent1  
 73 and parent2 are considered for the crossover process, and then combined in the form of ring, as shown in fig.  
 74 ?? decided in any point of ring. The children are created with a random number generated in any point of ring  
 75 according to the length of the combined two parental chromosomes. With reference to the cutting point, while  
 76 one of the children is created in the clockwise direction, the other one is created in direction of the anticlockwise,  
 77 as shown in fig. ??(c). Then swapping and reversing process is performed in the Ring Crossover operator, as  
 78 shown in fig. ??(d).

79 Figure ?? : Ring Crossover Procedure [11] The primary goals of this work are to produce a performance  
 80 comparison between traditional Brute force search algorithm and genetic algorithm with improved parameters  
 81 based method, and to determine the use of typical GA-based methods in the field of cryptanalysis.

82 The procedure to carry out the cryptanalysis using GA in order to break the key is as follows 1. Input:  
 83 ciphertext, and the language statistics. 2. Randomly generate an initial pool of solutions (keys).  $C_K = (i \oplus)$   
 84  $|K(i) \oplus u - D(i) \oplus (i, j \oplus \tilde{A}) | K(i, j) \oplus b - D(i, j) \oplus b | + (i, j, k \oplus \tilde{A}) | K(i, j, k) \oplus t - D(i, j, k) \oplus t | (!)$

85 Our objective in this paper is to compare the results obtained from Brute Force search algorithm with the  
 86 Genetic Algorithms with improved parameters. The experiments were conducted on Core chromosome, the more  
 87 times it is likely to be selected to reproduce. Crossover : Crossover selects genes from its parent chromosomes and  
 88 creates a new offspring. The simplest way to do this is to choose randomly some crossover point and everything  
 89 before this point is copied from the first parent and then, everything after a crossover point copied from the  
 90 second parent. Mutation : After a crossover, mutation is performed. This is to prevent falling all solutions in  
 91 population into a local optimum of solved problem. Mutation changes randomly the new offspring. In binary  
 92 GA we can switch a few randomly chosen bits from 1 to 0 or from 0 to 1. From the above table, it is found that  
 93 both GA works better than Brute force algorithm in terms of time taken as well as obtaining number of key bits.  
 94 In this paper, we have used a Genetic algorithm with Ring crossover and other operators for the cryptanalysis of  
 95 Simplified Data Encryption Standard. We found that Genetic Algorithm is far better than Brute Force search  
 96 algorithm for cryptanalysis of S-DES. Although S-DES is a simple encryption algorithm, GA with Ring Crossover  
 97 method can be adopted to handle other complex block ciphers like DES and AES. <sup>1 2</sup>

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Figure 1: Figure 1 :

2 I.

Figure 2: Figure 2 :

I

Figure 3: Global

3 INTRODUCTION

Figure 4: Figure 3 :

3 N

Figure 5: 3 .

2 II.

Figure 6: 2

5 R

Figure 7: Figure 5 :

6 RELATED

Figure 8: Figure 6 :

EI

Figure 9:

## GA Parameters

The following are the GA parameters used during the experimentation:

Population Size: 100

Selection : Tournament Selection operator

Crossover Ring Crossover

Crossover: .85

Mutation: .02

No. of Generation: 50

Genetic

Algorithm

Brute Force

Search

Algorithm

Figure 10: Comparison Between GA and Brute Force Search Algorithm

1

S. No	Amount of Cipher Text	Force Search Algorithm.		Time Taken by GA (M)	Time Taken by Brute Force search (M)
		No. of bits matched using GA	No. of bits matched using Brute Force search		
1.	200	5	5	4.7	24.3
2.	400	4	3	2.1	24.7
3.	600	7	6	1.9	23.6
4.	800	8	7	3.1	24.1
5.	1000	9	7	2.6	25.1
6.	1200	9	8	2.1	25.5

Figure 11: Table 1 :

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